

Gaussian Classical capacity of Gaussian Quantum channels

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Outline

- Quantum channels for information transmission
 - Classical capacity of quantum channels
- Gaussian channels
 - Additive (classical) noise quantum channel
- Gaussian capacity
- Mathematical challenges
- Conclusion (Gaussian)

Classical capacity of quantum channel

- Quantum channel – completely positive map $\rho^{out} = T[\rho^{in}]$
- Information encoding into quantum states – modulation

$\{P_i, \rho_i^{in}\}$ – input ensemble

- Average input: $\overline{\rho^{in}} = \sum_i P_i \rho_i^{in}$, average output $\overline{\rho} = \sum_i P_i \rho_i^{out}$
- Information content – Von Neumann entropy $S(\rho) = -Tr[\rho \cdot \text{Log}_2 \rho]$
- Information capacity [Holevo-Schumacher-Westmoreland]

$$C_1[T] = \max_{\{P_i, \rho_i^{in}\}} \left[S(\overline{\rho}) - \sum_i P_i S(\rho_i^{out}) \right]$$

Additivity problem

- When using entangled inputs
is classical capacity superadditive?

$$C_1(T_1 \otimes T_2) > C_1(T_1) + C_1(T_2)$$

– sometimes yes [Hastings]

- Capacity :
$$C(T) = \lim_{n \rightarrow \infty} \frac{1}{n} C_1(T^{\otimes n})$$

- Memory channels / correlated channel uses :

– Parallel n-mode channel

$$C(T) = \lim_{n \rightarrow \infty} \frac{1}{n} C_1(T^{(n)})$$

Bosonic channels – continuous variables

- Quadratures of electromagnetic field mode $\hbar = \omega = 1$

$$[\hat{a}, \hat{a}^\dagger] = 1 \quad \hat{q} = \frac{1}{\sqrt{2}}(\hat{a} + \hat{a}^\dagger) \quad \hat{p} = \frac{1}{i\sqrt{2}}(\hat{a} - \hat{a}^\dagger)$$

- *Gaussian states* are fully determined by $\rho_G \rightarrow (\vec{d}, V_{in})$

– displacement vector $\vec{d} = (q, p)^T$ – first moments

– covariance matrix (CM) V – second moments

- *Gaussian channels* transform first and second moments

$$\vec{d}_{out} = X\vec{d}_{in} + \vec{d} \quad , \quad X - \text{Real matrix}$$

$$V_{out} = XV_{in}X + Y \quad , \quad Y - \text{CM}$$

Additive noise channel

- Classical noise $X = I$ $V_{out} = V_{in} + Y$
 - “displacement” with a Gaussian probability distribution
- Input energy constraint $Tr[V_{in} + V_m] \leq E$
- Non-trivial optimization problem
- Challenge 1: Classical capacity of Gaussian channels is additive
 - True for “entanglement breaking” channels [Holevo 2008]
- Restriction to Gaussian states: Gaussian modulation
 - “displacement” with a Gaussian probability distribution
- Modulated output CM $\overline{V} = V_{in} + V_m + Y$

Gaussian capacity

- Von Neumann entropy is a function of CM

$$C_1[T_G] = \max_{\{Tr(V_{in} + V_m) \leq E\}} \left[S(\bar{V}) - \int dq dp P(q, p) S(V_{out}) \right]$$

- All signal states at the output have the same entropy

$$C_1[T_G] = \max_{\{Tr(V_{in} + V_m) \leq E\}} \left[S(\bar{V}) - S(V_{out}) \right]$$

- Challenge 2: Gaussian capacity is the capacity

- Proven for lossy channel with vacuum noise
[V. Giovannetti et al., Phys. Rev. Lett. 92, 027902 (2004)]

One mode problem

- 2x2 CM: Symplectic eigenvalue $V \rightarrow \text{diag}(v, v)$
- Entropy: $S(V) = g(v - 1/2)$

$$g(x) = (x+1)\text{Log}_2(x+1) - (x)\text{Log}_2(x)$$

- Optimization in terms of CM

$$C_1[T_G] = \max_{\{Tr(V_{in} + V_m) \leq E\}} [S(V_{in} + V_m + Y) - S(V_{in} + Y)]$$

- If noise CM is diagonalized by a passive symplectic transformation then optimal input and modulation are diagonal In the eigenbasis of the CM of noise

$$\bar{V} \rightarrow \bar{v} = \sqrt{(i_q + m_q + e_q)(i_p + m_p + e_p)}$$

Waterfilling solution

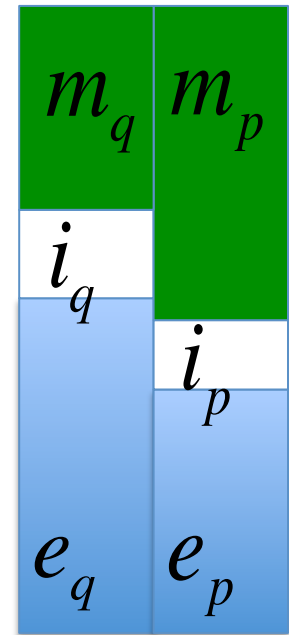
- Uniform distribution of the output energy is optimal
- Optimal input is a *minimum output entropy* Gaussian state

$$C_1[T] = g\left(\left(E + e_q + e_p - 1\right) / 2\right) - g\left(\sqrt{e_q e_p}\right) \quad \bar{V}_q = \bar{V}_p$$

- Challenge 3: minimum output entropy is achieved on a Gaussian state (vacuum state)

$$E > E_{th} = \sqrt{\frac{e_q}{e_p}} + e_q - e_p$$

$$\frac{\dot{i}_q}{\dot{i}_p} = \frac{e_q}{e_p}$$



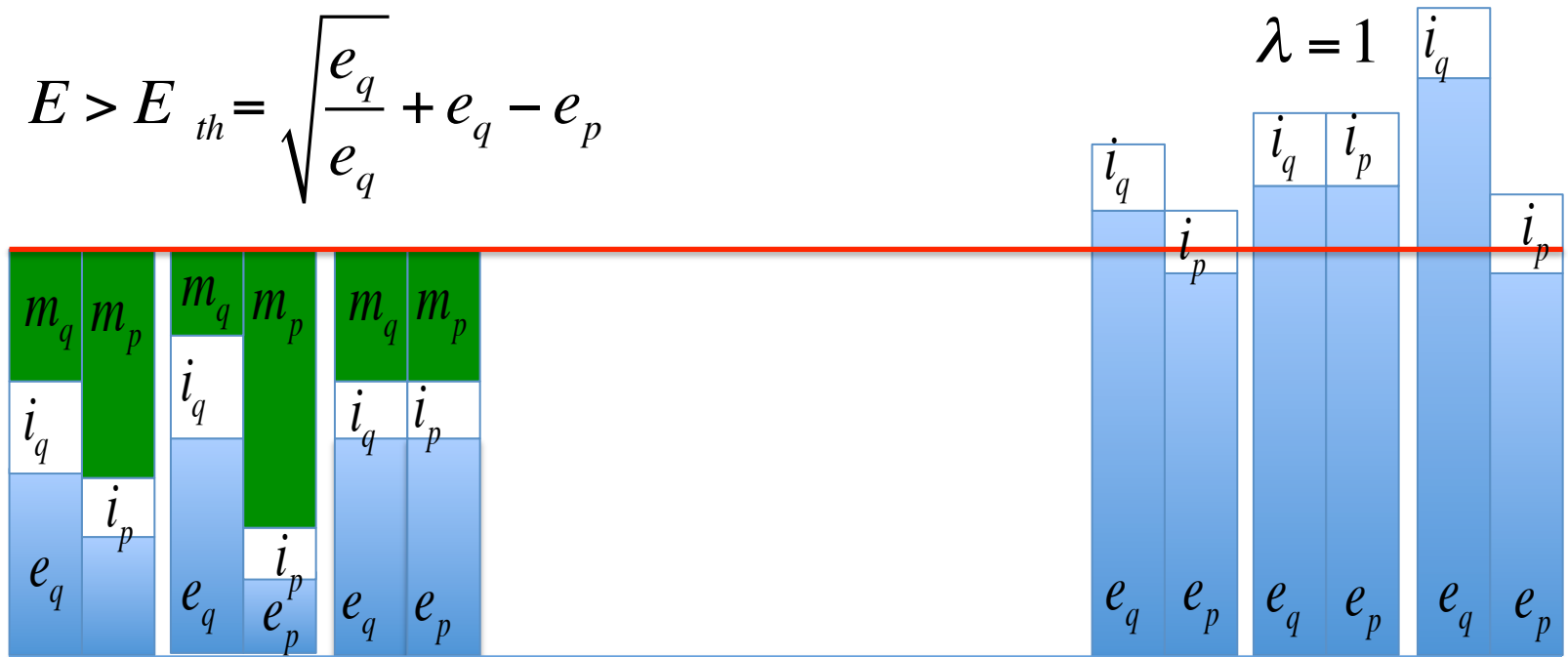
[Holevo, et al. PRA 59, 1820 (1999)]

Optimal states for multimode channel

Non-trivial input energy distribution between the modes

– common Lagrange multiplier – **waterfilling**

$$E > E_{th} = \sqrt{\frac{e_q}{e_p}} + e_q - e_p$$



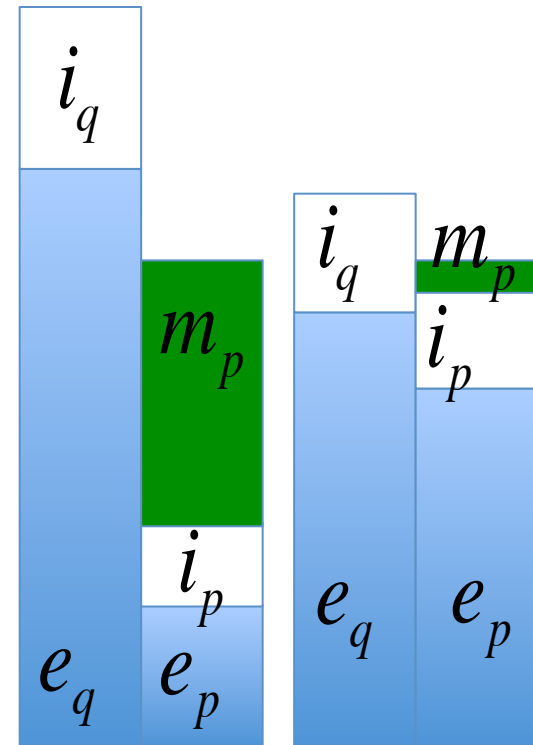
$$\frac{i_q}{i_p} = \frac{e_q}{e_p}$$

$$\frac{i_q}{i_p} = 1$$

$$1 < E < E_{th}$$

Solution of 2^d type

- Only one quadrature modulated
- Lagrange multipliers – no explicit solution
 - An implicit solution given by a transcendent equation
 - Input energy also spent on the non-modulated quadrature
 - Optimal input state is not the minimum output entropy state



$$1 < \frac{i_q}{i_p} < \frac{e_q}{e_p}$$

[Schäfer, et al. PRA 80, 062313 (2009)]

Solution of 2^d type

$$\frac{g'(\bar{v} - 1/2)}{\bar{v}}(E + e_p - e_q - 2i_q) = \frac{g'(v_{out} - 1/2)}{v_{out}} \left(e_p - e_q \frac{i_p}{i_q} \right)$$

$$g'(x) = \text{Log}_2(x+1) - \text{Log}_2 x$$

$$\bar{v} = \sqrt{(i_q + e_q)(E - i_q + e_p)}$$

$$v_{out} = \sqrt{(i_q + e_q)(i_p + e_q)}$$

$$i_q + i_p + m_p = E$$

$$i_q i_p = 1/4$$

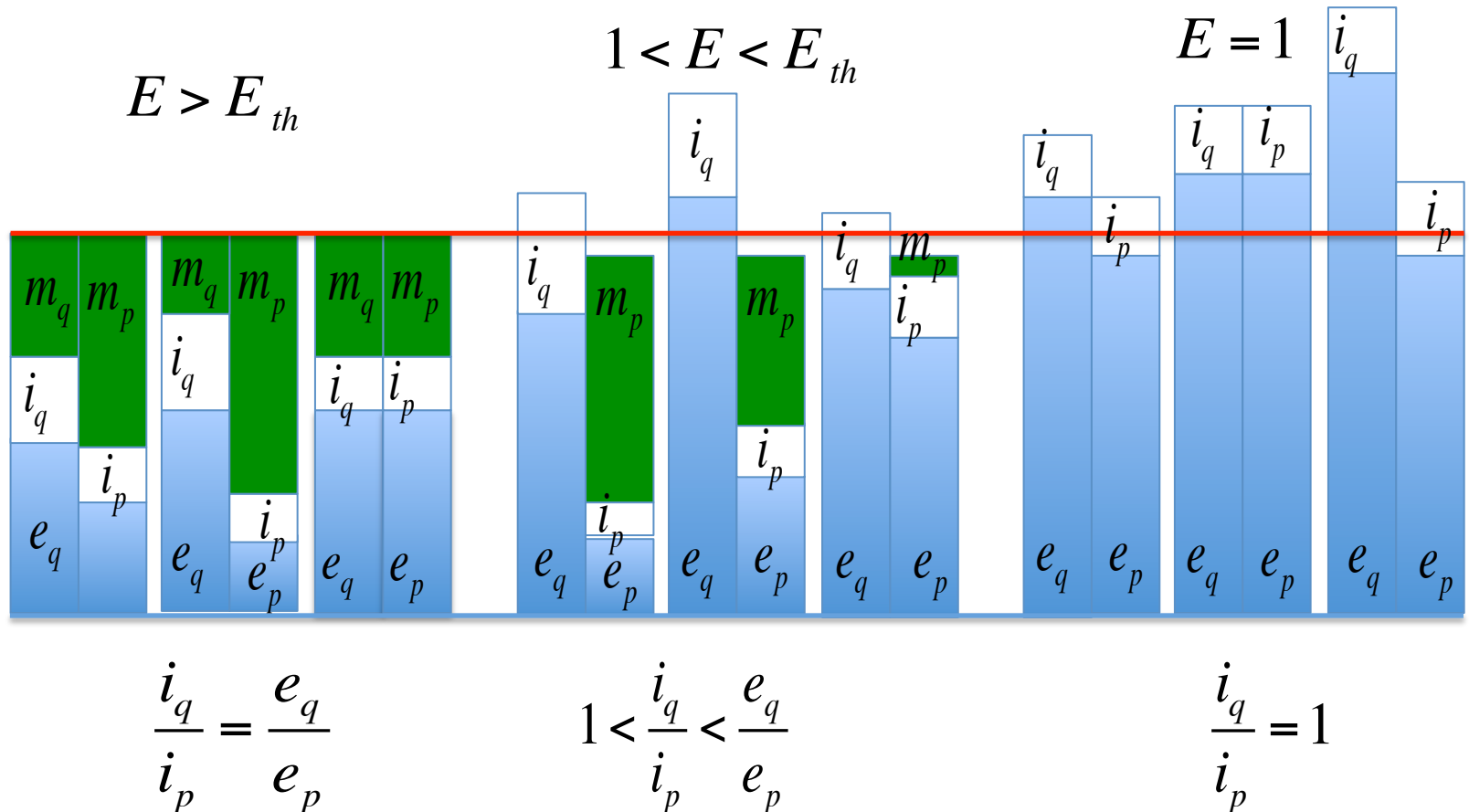
$$\frac{g'(\bar{v} - 1/2)}{\bar{v}}(i_q + e_q) = \lambda$$

- Lagrange multiplier

Optimal states for multimode channel

Non-trivial energy distribution among the modes

– common Lagrange multiplier – generalized **waterfilling**



Application - memory channel

- Correlated noise \rightarrow correlated optimal input
– multimode problem

$$C(T) = \lim_{n \rightarrow \infty} \frac{1}{n} C_1(T^{(n)})$$

- Parallel channel

- Gauss – Markov noise

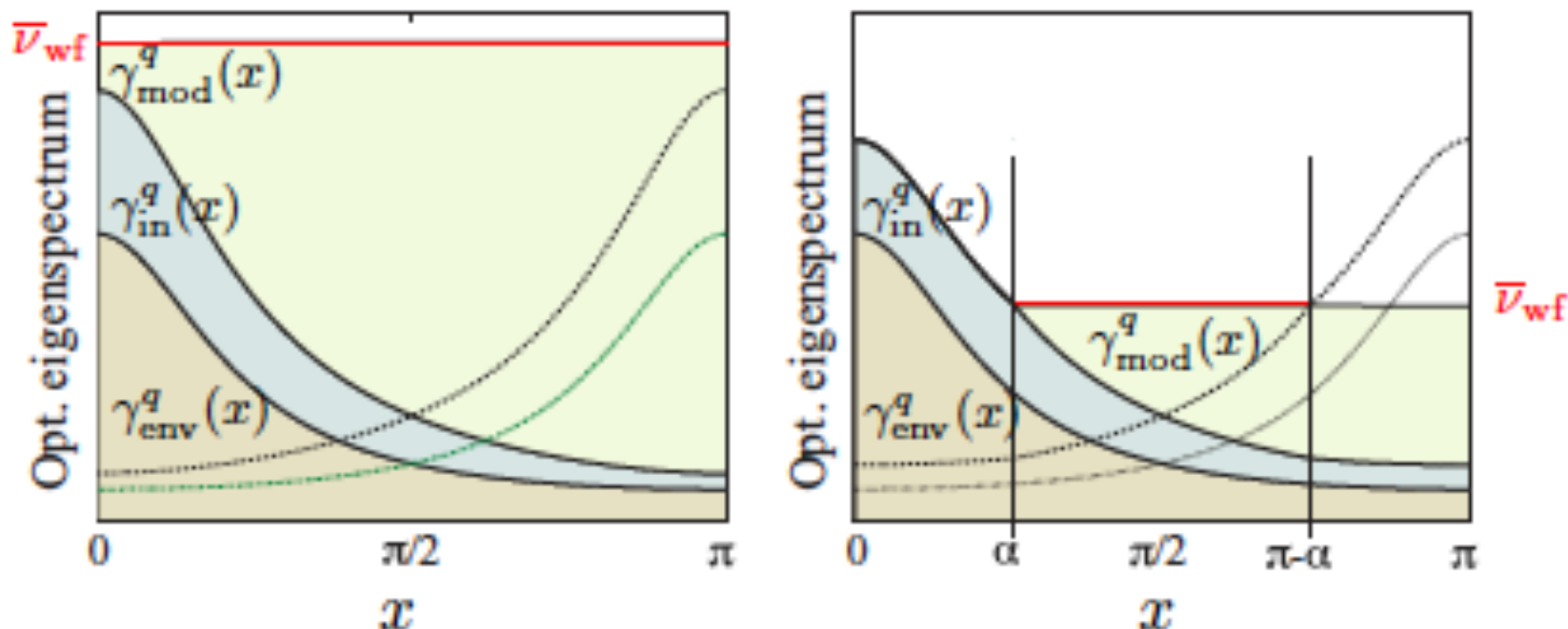
$$\{M(\phi)\}_{i,j} = N\phi^{|i-j|}$$

$$Y = \begin{pmatrix} M_q(\phi) & 0 \\ 0 & M_p(-\phi) \end{pmatrix}$$

- Symplectic and orthogonal diagonalization is possible as $M_q(\phi)$ and $M_q(-\phi)$ commute in the limit $n \rightarrow \infty$

Additive Gauss-Markov noise

Diagonalization of infinite noise matrix
– continuous spectrum on a finite domain



Optimal multimode input state in the original basis is entangled – **preparation?**

Conclusion

- Gaussian capacity of additive noise bosonic channels
- Three types of solution for one mode
- Input state that realizes classical capacity is not always the minimum output entropy state
- Challenge 1: Minimum output entropy for Gaussian channels is achieved on a Gaussian state
- Challenge 2: Classical capacity of memoryless Gaussian quantum channels is realized by Gaussian states
- Challenge 3: Classical capacity of memoryless Gaussian quantum channels is additive

Thank you for your attention !

